

# CSC 370 - Computer Organization and Architecture

## Lecture 0: Review of Numeric Representation

### Number Systems - Base 10

The number system that we use is base 10:

$$\begin{aligned}1734 &= 1000 + 700 + 30 + 4 \\ &= 1 \times 1000 + 7 \times 100 + 3 \times 10 + 4 \times 1 \\ &= 1 \times 10^3 + 7 \times 10^2 + 3 \times 10^1 + 4 \times 10^0\end{aligned}$$

$$\begin{aligned}724.5 &= 7 \times 100 + 2 \times 10 + 4 \times 1 + 5 \times 0.1 \\ &= 7 \times 10^2 + 2 \times 10^1 + 4 \times 10^0 + 5 \times 10^{-1}\end{aligned}$$

Why use base 10?

## Number Systems - Base 2

For computers, base 2 is more convenient (why?)

$$10011_2 = 1 \times 16 + 0 \times 8 + 0 \times 4 + 1 \times 2 + 1 \times 1 = 19_{10}$$

$$100010_2 = 1 \times 32 + 0 \times 16 + 0 \times 8 + 0 \times 4 + 1 \times 2 + 0 \times 1 = 34_{10}$$

$$\begin{aligned} 101.001_2 &= 1 \times 4 + 0 \times 2 + 1 \times 1 + 0 \times 0.5 + 0 \times 0.25 + 1 \times 0.125 \\ &= 5.125_{10} \end{aligned}$$

Example -  $1101011_2 = ?$

$$10110111_2 = ?$$

$$10100.1101_2 = ?$$

## Number Systems - Base 16

Hexadecimal (base 16) numbers are commonly used because it is convert them into binary (base 2) and vice versa.

$$\begin{aligned} 8CE_{16} &= 8 \times 256 + 12 \times 16 + 14 \times 1 \\ &= 2048 + 192 + 14 \\ &= 2254 \end{aligned}$$

$$\begin{aligned} 3F9 &= 3 \times 256 + 15 \times 16 + 9 \times 1 \\ &= 768 + 240 + 9 = 1017 \end{aligned}$$

## Number Systems - Base 16 (continued)

Base 2 is easily converted into base 16:

$$100011001110_2 = 1000 \ 1100 \ 1110 = 8 \ C \ E_{16}$$

$$11101101110101001_2 = 1 \ 1101 \ 1011 \ 1010 \ 1001 = 1 \ D \ B \ A \ 9_{16}$$

$$10110001010000010111_2 = ?_{16}$$

$$101101010010111011_2 = ?_{16}$$

## Number Systems - Base 16 (continued)

Converting base 16 into base 2 works the same way:

$$F3A5_{16} = 1111 \ 0011 \ 1010 \ 0101_2$$

$$76EF_{16} = 0111 \ 0110 \ 1110 \ 1111_2$$

$$AB3D_{16} = ?_2$$

$$15C.38_{16} = ?_2$$

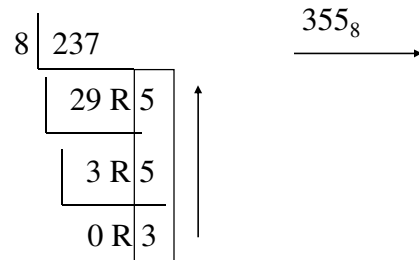
## Converting From Decimal to Binary

$$\begin{array}{r|l} 19 & \\ \hline 9 & \text{R } 1 \\ \hline 4 & \text{R } 1 \\ \hline 2 & \text{R } 0 \\ \hline 1 & \text{R } 0 \\ \hline 0 & \text{R } 1 \end{array} \quad \uparrow \quad 10011_2$$

## Converting From Decimal to Hexadecimal

$$16 \left| \begin{array}{l} 237 \\ \hline 14 \text{ R } 13 \\ \hline 0 \text{ R } 14 \end{array} \right. \quad \xrightarrow{\text{ED}_{16}} \quad \uparrow$$

## Converting From Decimal to Octal



## Binary, Octal, Decimal and Hexadecimal Equivalents

Binary	Decimal	Octal	Hex.	Binary	Decimal	Octal	Hex.
0000	0	0	0	1000	8	10	8
0001	1	1	1	1001	9	11	9
0010	2	2	2	1010	10	12	A
0011	3	3	3	1011	11	13	B
0100	4	4	4	1100	12	14	C
0101	5	5	5	1101	13	15	D
0110	6	6	6	1110	14	16	E
0111	7	7	7	1111	15	17	F

## 2s Complement Representation

- In 2s complement representation, we subtract the absolute value from  $2^n$ :

$$\begin{array}{r} 10000000 \\ \underline{0000110} \\ 11111010 \end{array}$$

$$\begin{array}{r} -6 \\ \underline{+ +13} \\ + 7 \end{array} \qquad 1 \qquad \begin{array}{r} 11111010 \\ \underline{00001101} \\ 00000111 \quad (= +7) \end{array}$$

## 2s Complement Representation (continued)

- The 2s complement representation can also be found by reversing the bits (into 1s complement) and then adding 1:

$$\begin{array}{r} 6 \Rightarrow 00000110 \Rightarrow 11111001 \\ \quad \quad \quad + \quad \quad \quad \underline{1} \\ \quad \quad \quad \quad \quad \quad 11111010 \end{array}$$

$$\begin{array}{r} 43 \Rightarrow 00101011 \Rightarrow 11010100 \\ \quad \quad \quad + \quad \quad \quad \underline{1} \\ \quad \quad \quad \quad \quad \quad 11010101 \end{array}$$

## Overflow

- If an addition operation produces a result that exceeds our number system's range, **overflow** has occurred.
- Addition of two numbers of the same sign produces overflow; addition two numbers of opposite sign cannot cause overflow.

$$\begin{array}{r}
 -3 \quad 1101 \\
 +6 \quad \underline{0110} \\
 +3 \quad 1\ 0011 = +3
 \end{array}
 \qquad
 \begin{array}{r}
 +5 \quad 0101 \\
 +6 \quad \underline{0110} \\
 +11 \quad 1011 = -5
 \end{array}$$

$$\begin{array}{r}
 -8 \quad 1000 \\
 -8 \quad \underline{1000} \\
 -16 \quad 1\ 0000 = 0
 \end{array}
 \qquad
 \begin{array}{r}
 +7 \quad 0111 \\
 +7 \quad \underline{0111} \\
 +14 \quad 1110 = -2
 \end{array}$$

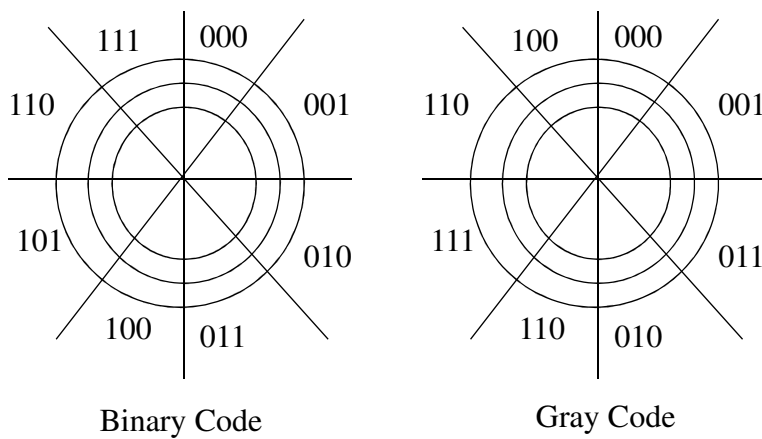
## Binary Representation of Decimal Numbers

<u>Decimal Digit</u>	<u>BCD (8421)</u>	<u>2421</u>	<u>Excess-3</u>
0	0000	0000	0011
1	0001	0001	0100
2	0010	0010	0101
3	0011	0011	0110
4	0100	0100	0111
5	0101	1011	1000
6	0110	1100	1001
7	0111	1101	1010
8	1000	1110	1011
9	1001	1111	1100

## Gray Codes

- Sometimes electromechanical applications of digital systems (machine tools, automotive brake systems and copiers) require a digital value that indicates a mechanical position.
- A standard binary code may see more than one bit change from one position to another, which could lead to an incorrect reading if mechanical assembly is imperfect.

## Binary Code vs. Gray Code





## ASCII representation of characters

- ASCII (*American Standard Code for Information Interchange*) is a numeric code used to represent characters.
- All characters are represented this way including:
  - words (character strings)
  - numbers
  - punctuation
  - control characters
- There are separate values for upper case and lower case characters:

A	65	z	122
B	66	<i>blank</i>	32
Z	90	\$	52
a	97	0	48
b	98	9	57

## Control Codes

- ASCII (a 7-bit code) has  $2^7 = 128$  values.
- We only need 62 for alphanumeric characters. Even after accounting for common punctuation, there are far more available code values than we need. What do we use them for?
- Control codes include DEL (for delete), NUL (for null), STX (Start of Text), CR (for carriage return), etc.

## Error Detection Codes

- An error is a corruption of the data from its correct state.
- There are several codes that allow use to detect an error. These include:
  - Parity
  - CRC
  - Checksum

## Parity

- Parity is an extra bit appended to our data which indicates whether the data bits add up to an even (for even parity) or odd (for odd parity) value.

## Parity Generation

<u>Message (xyz)</u>	<u>P(odd)</u>	<u>P(even)</u>
000	1	0
001	0	1
010	1	0
011	0	1
100	1	0
101	0	1
110	1	0
111	0	1

### Odd Parity

