

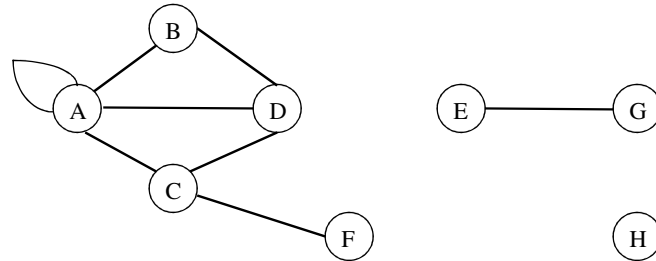
CSC 344 – Algorithms and Complexity

Lecture #12 – Graphs (Extended)

What is a Graph?

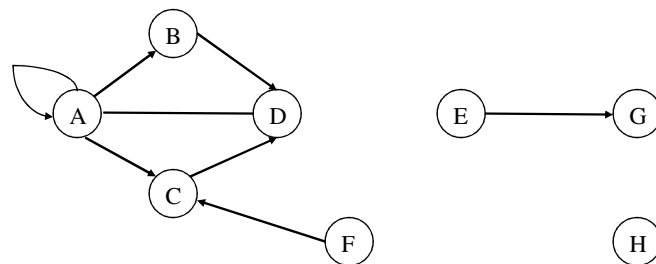
- A graph consists of a set of *nodes* (or *vertices*) and a set of *arcs* (or *edges*).
- Each arc in a graph is specified by a pair of nodes.
- If the pair of nodes that make up the arcs are *ordered pairs* then the graph is a *directed graph* or *digraph*.

Undirected Graph – An Example



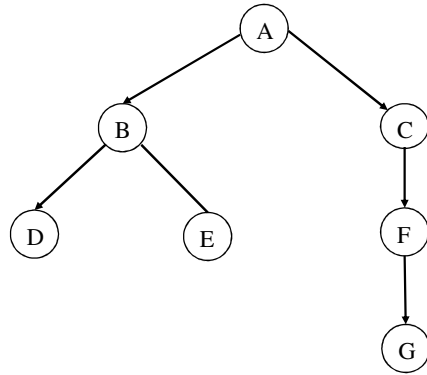
Set of nodes = {A, B, C, D, E, F, G, H}
Set of arcs = {(A, B), (A, D), (A, C), (C, D),
(C, F), (E, G), (A, A)}

Directed Graph – An Example

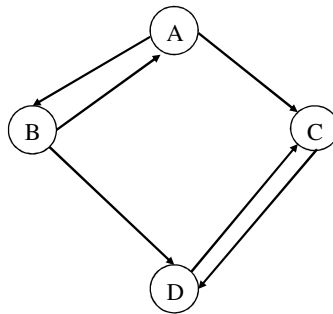


Set of nodes = {A, B, C, D, E, F, G, H}
Set of arcs = {<A, B>, <A, D>, <A, C>, <C, D>,
<F, C>, <E, G>, <A, A>}

Digraph – An Example



Digraph – An Example



A graph need not be a tree but a tree must be a graph.

Other Definitions

- A node n is incident to an arc x if n is one of the two nodes in the ordered pair of nodes constituting x . We also say that x is incident to n .
- The *degree of a node* is the number of arcs incident to it.
- *indegree of n* – the number of arcs with n as the head.
- *outdegree of n* – the number of arcs with n as the tail.

Weighted Graphs

- A number may be associated with each arc of a graph. Such a graph is called a *weighted graph* or *network*. The number associated with an arc is called the *weight*.

Operations Used With Graphs

- $join(a, b)$ – adds an arc from node a to b .
- $joinwt(a, b, x)$ – adds an arc from a to b with weight x .
- $remove(a, b)$ – removes an arc from a to b if it exists.
- $removewt(a, b, x)$ – removes an arc from a to b and sets x to the weight of the now-defunct arc.

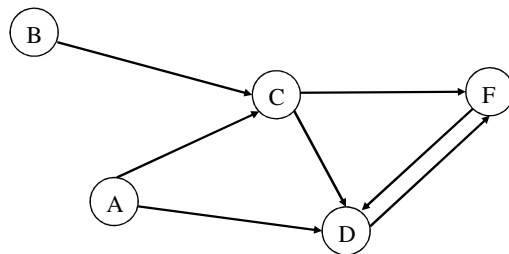
Paths and Cycles

- A path of length k from node a to node b is defined as a sequence of $k + 1$ nodes n_1, n_2, \dots, n_{k+1} such that $n_1 = a$ and $n_{k+1} = b$ and $adjacent(n_i, n_{i+1})$ is true for all i between 1 and k .
- A path from one node to itself is called a **cycle**.
- A graph with a cycle is cyclic; a graph without cycles is acyclic.
- Directed Acyclic Graphs are called ***dags***.

Transitive Closure

- Let's assume that the adjacency matrix (*adj*) completely describes the graph (the nodes contain no data and the graph is unweighted).
 - `if (adj[i][k] && adj[k][j] == true)`
 - `// we have a 2-arc path from i to j`
- What if the path requires 3 or more arcs?

Sample Graph



adj

	A	B	C	D	E
A	0	0	1	1	0
B	0	0	1	0	0
C	0	0	0	1	1
D	0	0	0	0	1
E	0	0	0	1	0

*adj*₂

	A	B	C	D	E
A	0	0	0	1	1
B	0	0	0	1	1
C	0	0	0	1	1
D	0	0	0	1	0
E	0	0	0	0	1

*adj*₃

	A	B	C	D	E
A	0	0	0	1	1
B	0	0	0	1	1
C	0	0	0	1	1
D	0	0	0	0	1
E	0	0	0	1	0

*adj*₄

	A	B	C	D	E
A	0	0	0	1	1
B	0	0	0	1	1
C	0	0	0	1	1
D	0	0	0	1	0
E	0	0	0	1	1

$$path = adj_1 | adj_2 | adj_3 | adj_4 | adj_5$$

	A	B	C	D	E
A	0	0	1	1	1
B	0	0	1	1	1
C	0	0	0	1	1
D	0	0	0	1	1
E	0	0	0	1	1

Graph.h

```
#ifndef __GRAPH__
#define __GRAPH__
#endif

using namespace std;

const int MaxNodes = 50;
typedef int NodeStuffType;

struct node {
    NodeStuffType data;
};

struct arc {
    bool adj;
};
```

```

class Graph
{
public:
    Graph(void);
    void  join(int node1, int node2);
    void  remove(int node1, int node2);
    bool  adjacent(int node1, int node2);
    void  transClose(int path[][MaxNodes]);
private:
    void  prod(int a[][MaxNodes],
              int c[][MaxNodes]);
    struct node nodes[MaxNodes];
    struct arc  arcs[MaxNodes][MaxNodes];
};

```

Graph.cpp

```

#include "Graph.h"

Graph::Graph(void)
{
    int i, j;
    for (i = 0; i < MaxNodes; i++)
        for (j = 0; j < MaxNodes; j++)
            arcs[i][j].adj = false;
}

void  Graph::join(int node1, int node2)  {
    arcs[node1][node2].adj = true;
}

```

```

void Graph::remove(int node1, int node2) {
    arcs[node1][node2].adj = false;
}
bool Graph::adjacent(int node1, int node2)    {
    return((arcs[node1][node2].adj == true)?
           true : false);
}

```

```

void Graph::transClose(int path[][MaxNodes])    {
    int i, j, k;
    int newprod[MaxNodes][MaxNodes],
        adjprod[MaxNodes][MaxNodes];

    for (i = 0; i < MaxNodes; i++)
        for (j = 0; j < MaxNodes; j++)
            adjprod[i][j] = path[i][j]
                = arcs[i][j].adj;

    for (i = 1; i < MaxNodes; i++)    {
        // i represents the number of times adj
        // has been mulitplied by itself to
        // obtain adjprod.  At this point path
        // represents all paths of length i or
        // less
    }
}

```

```

    prod(adjprod, newprod);
    for (j = 0; j < MaxNodes; j++)
        for (k = 0; k < MaxNodes; k++)
            path[j][k]
                = path[j][k] || newprod[j][k];

    for (j = 0; j < MaxNodes; j++)
        for (k = 0; k < MaxNodes; k++)
            adjprod[j][k] = newprod[j][k];
    }
}

```

```

void Graph::prod(int a[][MaxNodes],
                int c[][MaxNodes]) {
    int i, j, k, val;

    for (i = 0; i < MaxNodes; i++)
        //pass through rows
        for (j = 0; j < MaxNodes; j++) {
            // pass through columns
            val = false;
            for (k = 0; k < MaxNodes; k++)
                val
                    = val ||
                      (a[i][k] && arcs[i][j].adj);
            c[i][j] = val;
        } // for j..
    }

```

Shortcoming in **transClosure ()**

- The matrix multiplication that is performed is $O(n^3)$. It is performed $n-1$ times. That makes the efficiency of the algorithm $O(n^4)$, which is generally unacceptable.

Warshall's Algorithm

- We need a more efficient algorithm.
- Matrix $path_k$ is defined such that $path_k[i][j]$ is true if and only if there is path from node i to j that does not pass through any node numbered higher than k .
- Can we determine $path_{k+1}$ from $path_k$?

Warshall's Algorithm

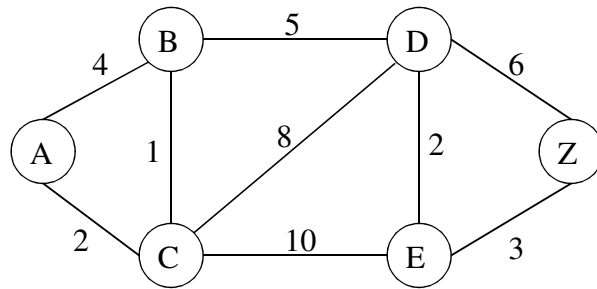
- $path_{k+1}$ will be true if and only if:
 1. $path_k[i][j] == true$
 2. $path_k[i][k+1] == true$
&& $path_k[k+1][j] == true$

quickTransClose ()

```
void Graph::quickTransClose
    (int path[][MaxNodes]) {
    int i, j, k;
    for (i = 0; i < MaxNodes; i++)
        for (j = 0; j < MaxNodes; j++)
            // Path starts off as adj
            path[i][j] = arcs[i][j].adj;

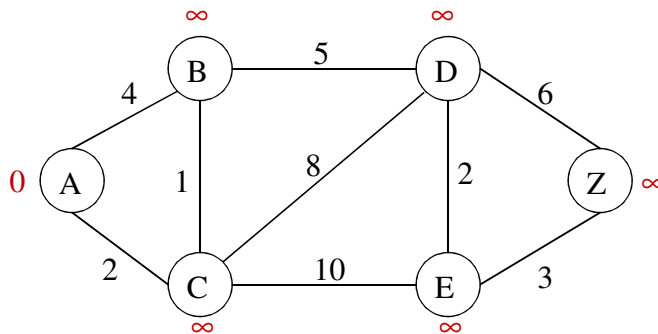
    for (k = 0; k < MaxNodes; k++)
        for (i = 0; i < MaxNodes; i++)
            if (path[i][k] == 1)
                for (j = 0; j < MaxNodes; j++)
                    path[i][j] = path[i][j] || path[k][j];
}
```

Dijkstra's Algorithm



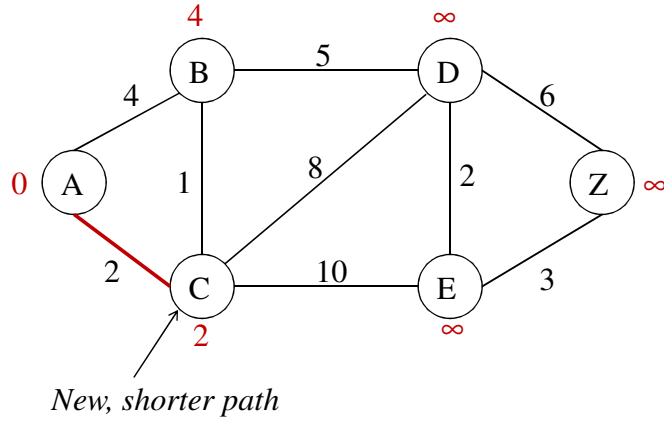
We want to find the shortest path from A to Z

Dijkstra's Algorithm



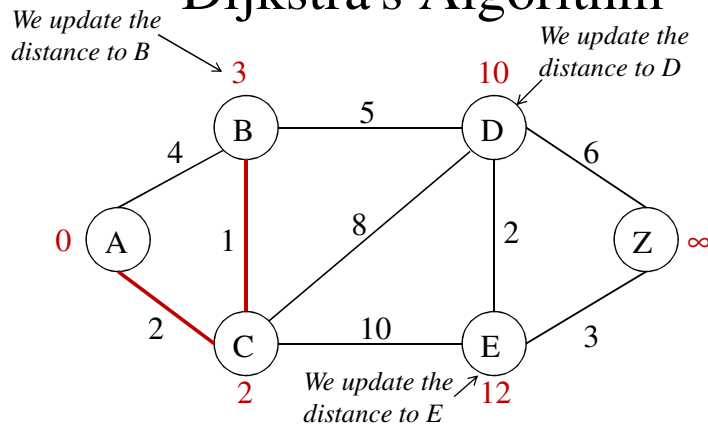
A's distance is 0; everyone else is initialize to ∞

Dijkstra's Algorithm



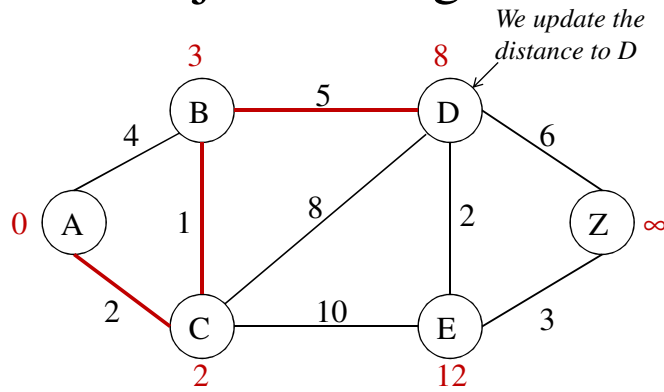
We update A's neighbor's

Dijkstra's Algorithm



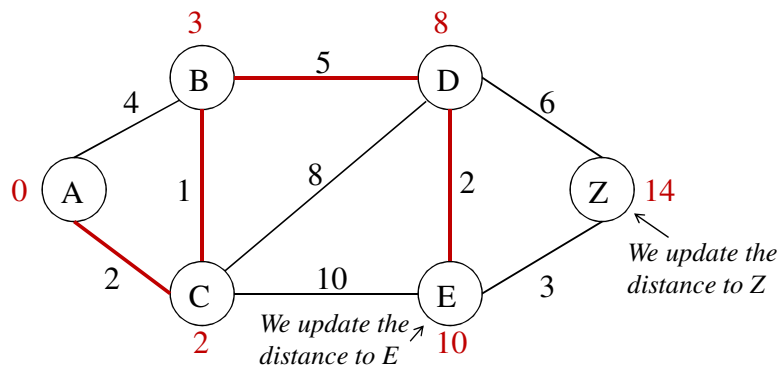
We update C's neighbor's since it is closer to A

Dijkstra's Algorithm



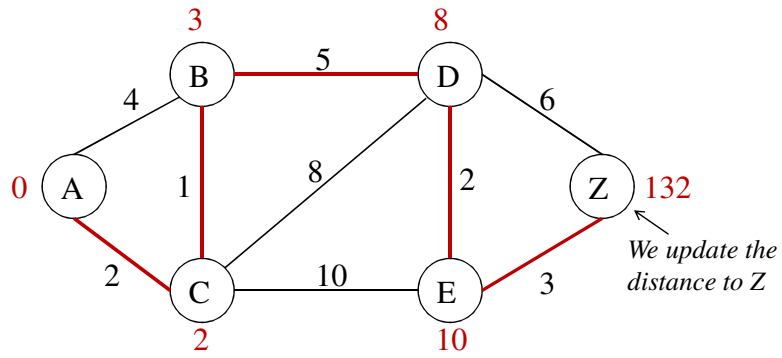
We update C's neighbor's since it is closer to A

Dijkstra's Algorithm



We update C's neighbor's since it is closer to A

Dijkstra's Algorithm



We update C's neighbor's
since it is closer to A