# Software II: Principles of Programming Languages

Lecture 7 – Expressions and Assignment Statements

# Why Expressions?

- Expressions are the fundamental means of specifying computations in a programming language
- To understand expression evaluation, need to be familiar with the orders of operator and operand evaluation
- Essence of imperative languages is dominant role of assignment statements

# Arithmetic Expressions

- Arithmetic evaluation was one of the motivations for the development of the first programming languages
- Arithmetic expressions consist of operators, operands, parentheses, and function calls

## Arithmetic Expressions: Design Issues

- Design issues for arithmetic expressions
  - Operator precedence rules?
  - Operator associativity rules?
  - Order of operand evaluation?
  - Operand evaluation side effects?
  - Operator overloading?
  - Type mixing in expressions?

# Arithmetic Expressions: Operators

- A unary operator has one operand
- A binary operator has two operands
- A ternary operator has three operands

# Arithmetic Expressions: Operator Precedence Rules

- The operator precedence rules for expression evaluation define the order in which "adjacent" operators of different precedence levels are evaluated
- Typical precedence levels
  - parentheses
  - unary operators
  - \*\* (if the language supports it)
  - **–** \*, /
  - +, -

# Arithmetic Expressions: Operator Associativity Rule

- The operator associativity rules for expression evaluation define the order in which adjacent operators with the same precedence level are evaluated
- Typical associativity rules
  - Left to right, except \*\*, which is right to left
  - Sometimes unary operators associate right to left (e.g., in FORTRAN)
- APL is different; all operators have equal precedence and all operators associate right to left
- Precedence and associativity rules can be overriden with parentheses

## Expressions in Ruby and Scheme

- Ruby
  - All arithmetic, relational, and assignment operators, as well as array indexing, shifts, and bitwise logic operators, are implemented as methods
  - One result of this is that these operators can all be overridden by application programs
- Scheme (and Common LISP)
  - All arithmetic and logic operations are by explicitly called subprograms
  - -a + b \* c is coded as (+ a (\* b c))

# Arithmetic Expressions: Conditional Expressions

- Conditional Expressions
- Appears in C-based languages (e.g., C, C++)
- An example:

```
average = (count == 0)? 0 : sum / count
```

• Evaluates as if written as follows:

```
if (count == 0)
  average = 0
else
  average = sum /count
```

# Arithmetic Expressions: Operand Evaluation Order

- Operand evaluation order
  - 1. Variables: fetch the value from memory
  - 2. Constants: sometimes a fetch from memory; sometimes the constant is in the machine language instruction
  - 3. Parenthesized expressions: evaluate all operands and operators first
  - 4. The most interesting case is when an operand is a function call

# Arithmetic Expressions: Potentials for Side Effects

- Functional side effects: when a function changes a two-way parameter or a non-local variable
- Problem with functional side effects:
- When a function referenced in an expression alters another operand of the expression; e.g., for a parameter change:

```
a = 10;
/* assume that fun changes its
parameter */
b = a + fun(&a);
```

#### **Functional Side Effects**

- 1. Two possible solutions to the problem
  - Write the language definition to disallow functional side effects
  - No two-way parameters in functions
  - No non-local references in functions
  - Advantage: it works!
  - <u>Disadvantage</u>: inflexibility of one-way parameters and lack of non-local references
- 2. Write the language definition to demand that operand evaluation order be fixed
  - **<u>Disadvantage</u>**: limits some compiler optimizations
  - Java requires that operands appear to be evaluated in left-toright order

### Referential Transparency

• A program has the property of *referential transparency* if any two expressions in the program that have the same value can be substituted for one another anywhere in the program, without affecting the action of the program

```
result1 = (fun(a) + b) / (fun(a) - c);
temp = fun(a);
result2 = (temp + b) / (temp - c);
```

- If fun has no side effects, result1 = result2
- Otherwise, not, and referential transparency is violated

### Referential Transparency (continued)

- Advantage of referential transparency
  - Semantics of a program is much easier to understand if it has referential transparency
- Because they do not have variables, programs in pure functional languages are referentially transparent
  - Functions cannot have state, which would be stored in local variables
  - If a function uses an outside value, it must be a constant (there are no variables). So, the value of a function depends only on its parameters

## Overloaded Operators

- Use of an operator for more than one purpose is called operator overloading
- Some are common (e.g., + for int and float)
- Some are potential trouble (e.g., \* in C and C++)
  - Loss of compiler error detection (omission of an operand should be a detectable error)
  - Some loss of readability

# Overloaded Operators (continued)

- C++, C#, and F# allow user-defined overloaded operators
  - When sensibly used, such operators can be an aid to readability (avoid method calls, expressions appear natural)
  - Potential problems:
    - Users can define nonsense operations
    - Readability may suffer, even when the operators make sense

## **Type Conversions**

- A <u>narrowing conversion</u> is one that converts an object to a type that cannot include all of the values of the original type e.g., **float** to **int**
- A <u>widening conversion</u> is one in which an object is converted to a type that can include at least approximations to all of the values of the original type e.g., int to float

# Type Conversions: Mixed Mode

- A <u>mixed-mode expression</u> is one that has operands of different types
- A *coercion* is an implicit type conversion
- Disadvantage of coercions:
  - They decrease in the type error detection ability of the compiler
- In most languages, all numeric types are coerced in expressions, using widening conversions
- In Ada, there are virtually no coercions in expressions
- In ML and F#, there are no coercions in expressions

# **Explicit Type Conversions**

- Called casting in C-based languages
- Examples
  - -C: (int)angle
  - -F#: float (sum)
- Note that F#'s syntax is similar to that of function calls

# **Errors in Expressions**

- Causes
  - Inherent limitations of arithmetic e.g., division by zero
  - Limitations of computer arithmetic e.g. overflow
- Often ignored by the run-time system

# **Relational Operators**

- Use relational operators and operands of various types
- Evaluate to some Boolean representation
- Operator symbols used vary somewhat among languages (!=, /=, ~=, .NE., <>, #)

### Other Relational Operators

- JavaScript and PHP have two additional relational operator, === and !==
- Similar to their cousins, == and !=, except that they do not coerce their operands
- Ruby uses == for equality relation operator that uses coercions and eql? for those that do not

# **Boolean Expressions**

- Operands are Boolean and the result is Boolean
- Example operators (&& || ! AND OR NOT)
- C89 has no Boolean type--it uses **int** type with **0** for false and nonzero for true
- One odd characteristic of C's expressions:
  - a < b < c is a legal expression, but the result is not what you might expect:
  - Left operator is evaluated, producing 0 or 1
  - The evaluation result is then compared with the third operand (i.e., c)

#### **Short Circuit Evaluation**

- An expression in which the result is determined without evaluating all of the operands and/or operators
- Example:

$$(13 * a) * (b / 13 - 1)$$

• If a is zero, there is no need to evaluate

$$(b /13 - 1)$$

#### **Short Circuit Evaluation**

• Problem with non-short-circuit evaluation

```
index = 0;
while (index <= length)
    && (LIST[index] != value)
    index++;</pre>
```

When index=length, LIST[index] will cause an indexing problem (assuming LIST is length - 1 long)

### Short Circuit Evaluation (continued)

- C, C++, and Java: use short-circuit evaluation for the usual Boolean operators (&& and ||), but also provide bitwise Boolean operators that are not short circuit (& and |)
- All logic operators in Ruby, Perl, ML, F#, and Python are short-circuit evaluated
- Ada: programmer can specify either (short-circuit is specified with and then and or else)

### Short Circuit Evaluation (continued)

- Short-circuit evaluation exposes the potential problem of side effects in expressions
- Examples

```
(a > b) || (b++ / 3)
(c != 0 && (c = getchar()) != '\n')
```

# **Assignment Statements**

• The general syntax

```
<target_var> <assign_operator> <expression>
```

- The assignment operator
  - = Fortran, BASIC, the C-based languages
  - := Ada, Pascal
  - = can be bad when it is overloaded for the relational operator for equality (that's why the C-based languages use == as the relational operator)

# Assignment Statements: Conditional Targets

```
Conditional targets (Perl)
($flag ? $total : $subtotal) = 0
Which is equivalent to
if ($flag) {
    $total = 0
} else {
    $subtotal = 0
}
```

# Assignment Statements: Compound Assignment Operators

- A shorthand method of specifying a commonly needed form of assignment
- Introduced in ALGOL; adopted by C and the C-based languages
- Example

```
a = a + b
```

• can be written as

a += b

# Assignment Statements: Unary Assignment Operators

- Unary assignment operators in C-based languages combine increment and decrement operations with assignment
- Examples

### Assignment as an Expression

• In the C-based languages, Perl, and JavaScript, the assignment statement produces a result and can be used as an operand

```
while ((ch = getchar())!= EOF) {...}
ch = getchar() is carried out; the result (assigned
to ch) is used as a conditional value for the while
statement
```

• Disadvantage: another kind of expression side effect

# Multiple Assignments

 Perl, Ruby, and Lua allow multiple-target multiple-source assignments

```
(\$first, \$second, \$third) = (20, 30, 40);
```

• Also, the following is legal and performs an interchange:

```
($first, $second) = ($second, $first);
```

### Assignment in Functional Languages

- Identifiers in functional languages are only names of values
- ML
  - Names are bound to values with val

```
val fruit = apples + oranges;
```

- If another val for fruit follows, it is a new and different name
- F#
  - F#'s let is like ML's val, except let also creates a new scope

# Mixed-Mode Assignment

- Assignment statements can also be mixed-mode
- In Fortran, C, Perl, and C++, any numeric type value can be assigned to any numeric type variable
- In Java and C#, only widening assignment coercions are done
- In Ada, there is no assignment coercion