CSC 160 Computer Programming for Non-Majors Day 16 (June 20, 2005)

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Review

- New syntax rule allows "local" definitions
- Can use for both variables and functions
- Common applications:
 - save recursive results to be used several times; improve *efficiency*
 - give names to intermediate results; improve *readability*
 - hide things "outside world" doesn't need to know about; improve *modularization*

Modularization

- Large project w/several programmers
- Each in charge of a "module" of the project
- For example, video game
 - one module in charge of rules of game
 - another module in charge of graphics, menus, buttons, etc.

"Open" scenario

- Jim changes one of his auxiliary functions to return a different type
- Julia, who was using Jim's auxiliary function in her module, finds her module mysteriously stop working
- Problem is hard to track down because Julia doesn't even know Jim has changed the function

Information-hiding scenario

- Each module has
 - "public interface" (like a function contract & examples)
 - "private implementation" (like a function body)
- Each programmer knows only the "public interface" of other programmers' modules
- Jim's auxiliary function isn't in his public interface, so nobody else is using it, so he can change it w/o screwing up Julia's module
- If Jim changes a publicly-known function, but it still satisfies its public interface, it won't screw up Julia's module

Interface vs. implementation

- Interface: "What you need to know in order to *use* the function(s)" How functions can be called & what they should return.
- Implementation: "What you need to know in order to *write* or *fix* the function(s)" How functions compute the result

Interface vs. implementation: Automobile analogy

- interface is steering wheel, pedals, gear shift...
- implementation is engine, carburetor, spark plugs, ...
- If repair shop changes positions of pedals, you have to learn how to drive again.
- If repair shop replaces spark plugs w/better model, but steering wheel, pedals, etc. still work as before, you don't have to change anything.

Review: operating on lists

```
; remove>10 : list-of-nums -> list-of-nums
(define (remove>10 nums)
(cond [(empty? nums) empty]
[(cons? nums)
(cond [(> (first nums) 10) (remove>10 (rest nums))]
[else (cons (first nums) (remove>10 (rest nums)))])))
```

```
"Examples of remove>10:"
(remove>10 empty) "should be" empty
(remove>10 (list 6)) "should be" (list 6)
(remove>10 (list 11)) "should be" empty
(remove>10 (list 6 11 10 -24 13 9)) "should be" (list 6 10 -24 9)
(remove>10 (list 11 10 -24 13 9)) "should be" (list 10 -24 9)
```

Review: generalizing the function

: remove>5 : list-of-nums -> list-of-nums ; remove>17: list-of-nums -> list-of-nums What these have in common is that they **remove all elements of the list greater** than a fixed threshold. So we generalize the function: ; remove-over: num list-of-nums -> list-of-nums (define (remove-over threshold nums) (cond [(empty? nums) empty] [(cons? nums) (cond [(> (first nums) threshold) (remove-over threshold (rest nums))] [else (cons (first nums) (remove-over threshold (rest nums)))])])) "Examples of remove-over:" (remove-over 6 empty) "should be" empty . . .

(remove-over 3.5 (list 4 9 17 2 6 3)) "should be" (list 2 3)

Generalizing the function *farther*

; remove<5 : list-of-nums -> list-of-nums ; remove>=4: list-of-nums -> list-of-nums ; remove-evens : list-of-nums -> list-of-nums

What all of these have in common is that they **perform a test on each element of the list, and remove the ones that pass the test.**

Generalization:

; remove-if : test list-of-nums -> list-of-nums

Q: What is a "test"?

A: a property that every number either has or doesn't have

A: a function from number to boolean

Note: change languages to Intermediate Student

Writing **remove-if**

; remove-if : (num -> boolean) list-of-nums -> list-of-nums (define (remove-if test? nums)

. . .

```
)

"Examples of remove-if:"

(remove-if even? (list 1 2 3 4 5)) "should be" (list 1 3 5)

(define (over-10? x) (> x 10))

(remove-if over-10? (list 3 12 10 5 16 -24 6)) "should be" (list 3 10 5 -24 6)

(define (under-5? x) (< x 5))

(remove-if under-5? (list 3 12 10 5 16 -24 6)) "should be" (list 12 10 5 16 6)
```

The routine stuff

```
; remove-if : (num -> boolean) list-of-nums -> list-of-nums
(define (remove-if test? nums)
 (cond [(empty? nums) empty]
       [(cons? nums)
        (cond [...
               (remove-if test? (rest nums))]
              [else
               (cons (first nums) (remove-if test? (rest nums)))])]))
"Examples of remove-if:"
(remove-if even? (list 1 2 3 4 5)) "should be" (list 1 3 5)
(define (over-10? x) (> x 10))
(remove-if over-10? (list 3 12 10 5 16 -24 6)) "should be" (list 3 10 5 -24 6)
(define (under-5? x) (< x 5))
(remove-if under-5? (list 3 12 10 5 16 -24 6)) "should be" (list 12 10 5 16 6)
```

Using the test

```
; remove-if : (num -> boolean) list-of-nums -> list-of-nums
(define (remove-if test? nums)
 (cond [(empty? nums) empty]
       [(cons? nums)
        (cond [(test? (first nums))
               (remove-if test? (rest nums))]
              [else
               (cons (first nums) (remove-if test? (rest nums)))])]))
"Examples of remove-if:"
(remove-if even? (list 1 2 3 4 5)) "should be" (list 1 3 5)
(define (over-10? x) (> x 10))
(remove-if over-10? (list 3 12 10 5 16 -24 6)) "should be" (list 3 10 5 -24 6)
(define (under-5? x) (< x 5))
(remove-if under-5? (list 3 12 10 5 16 -24 6)) "should be" (list 12 10 5 16 6)
```

Writing functions using remove-if

; remove<5 : list-of-nums -> list-of-nums
(define (under-5? x) (< x 5))
(define (remove<5 nums) (remove-if under-5? nums))</pre>

; remove>=7: list-of-nums -> list-of-nums You try this one.

; remove-evens : list-of-nums -> list-of-nums (define (remove-evens nums) (remove-if even? nums))

Another example

```
"Examples of cube-each:"
(cube-each empty) "should be" empty
(cube-each (list 2)) "should be" (list 8)
(cube-each (list 3 -2 0 5 -6)) "should be"
(list 27 -8 0 125 -216)
```

Similar functions

- ; sqrt-each : list-of-nums -> list-of-nums
- ; negate-each : list-of-nums -> list-of-nums
- What these have in common is that they **do the same thing to** each element of a list, returning a list of the results.
- So we generalize the functions:
- ; do-to-each : operation list-of-nums -> list-of-nums
- What's an "operation"? In this case, a function from number to number.
- ; do-to-each : (num -> num) list-of-nums -> list-of-nums

Writing do-to-each

```
; do-to-each : (num -> num) list-of-nums -> list-of-nums
(define (do-to-each op nums)
      (cond [(empty? nums) empty]
                    [(cons? nums)
                         (cons (op (first nums))
                               (do-to-each op (rest nums)))]))
"Examples of do-to-each:"
(do-to-each cube (list 3 5 -2)) "should be" (list 27 125 -8)
(do-to-each sqrt (list 4 25 0)) "should be" (list 2 5 0)
(do-to-each - (list 3 -2 0 7.5)) "should be" (list -3 2 0 -7.5)
```

Writing functions using do-to-each

; sqrt-each : list-of-nums -> list-of-nums (define (sqrt-each nums) (do-to-each sqrt nums))

; add-3-to-each : list-of-nums -> list-of-nums (define (add3 x) (+ x 3)) (define (add-3-to-each nums) (do-to-each add3 nums))

Generalizing the contract

Nothing in **remove-if** or **do-to-each** actually depends on *numbers*

Real contracts are

; remove-if : (X -> boolean) list-of-X -> listof-X

; do-to-each : (X -> X) list-of-X -> list-of-X where X is *any* type

Writing functions using these

; fire-over-100K : list-of-emps -> list-of-emps ; Auxiliary function earns-over-100K? : emp -> boolean (define (earns-over-100K? emp) (> (emp-salary emp) 100000)) (define (fire-over-100K emps) (remove-if earns-over-100K? emps))

; give-10%-raises: list-of-emps -> list-of-emps ; Auxiliary function give-10%-raise : emp -> emp (define (give-10%-raise emp) (make-emp (emp-name emp) (emp-id emp) (* 1.1 (emp-salary emp)))) (define (give-10%-raises emps) (do-to-each give-10%-raise emps))

Generalizing even farther

Nothing in **do-to-each** requires input and output lists to be the *same* type

Real contract is

- ; do-to-each : (X -> Y) list-of-X -> list-of-Y
- where X and Y are *any two* types, possibly the same.

Writing functions using this

; extract-names : list-of-emps -> list-of-strings (define (extract-names emps) (do-to-each emp-name emps))

"Example of extract-names:"

(extract-names (list (make-emp "Joe" 1 75000) (make-emp "Mary" 2 79995) (make-emp "Phil" 3 26000))) "should be" (list "Joe" "Mary" "Phil")

Dumb single-use functions

; add-3-to-each : list-of-nums -> list-of-nums (define (add3 x) (+ x 3)) (define (add-3-to-each nums) (do-to-each add3 nums))

Better: hide **add3** inside a local definition (define (add-3-to-each nums) (local [(define (add3 x) (+ x 3))] (do-to-each add3 nums)))

Could do the same thing with earns-over-100K? and give-10%-raise

An example where we *have* to use **local**

Note: we *couldn't* have defined **over-threshold?** outside **remove-over**, because it would have depended on the threshold value.

- ; add-up : list-of-nums -> num
- ; multiply-all : list-of-nums -> num
- ; largest : non-empty-list-of-nums -> num
- ; highest-paid : non-empty-list-of-emps -> emp
- What these have in common is that they *combine pairs* of objects to get a third object, repeatedly until whole list has been combined
- So we generalize. Note that in each case, we need to know what value to start with...

; combine : X (X X -> X) list-of-X -> X
(define (combine start-value combiner values)
...)

"Examples of combine:"

(define (add-up nums) (combine 0 + nums))

; insert standard test cases for add-up here

```
(define (multiply-all nums)
(combine 1 * nums))
```

; insert standard test cases for multiply-all here

(define (largest nums)
 (local [(define (larger num1 num2)
 (cond [(> num1 num2) num1]
 [else num2]))]
 (combine (first nums) larger (rest nums))))
; insert standard test cases for largest here

- In fact, there's no rule that the types of list elements and the type of the result are the same...
- ; combine : Y (X Y -> Y) list-of-X -> Y

For example,

Defining functions without names

(+3(*45))

doesn't require defining a variable to hold the value of (* 4 5), and then adding 3 to it!

Why should **add-3-to-each** require defining a function to add 3 to things, and then applying **do-to-each** to it?

Note: change languages to Intermediate Student with Lambda

Defining functions without names

New syntax rule: (lambda (param param ...) expr) constructs a function without a name and returns it. Example:

(define (add-3-to-each nums) (do-to-each (lambda (x) (+ x 3)) nums))

Defining functions without names

- Anything you can do with **lambda** can also be done with **local**; may be more readable because things have names
- Anything you can do with **local** can also be done with **lambda**, often a little shorter